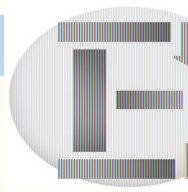


Run-time Checking/Monitoring at SDL Fredhopper

Stijn de Gouw

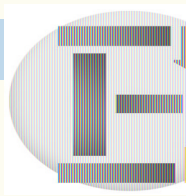




SDL Fredhopper

e-Commerce company, back-end for 300+ online shops





Communication in Java

Messages between Objects (sender/receiver):

Method Calls

▮ `y.m(p)` Current object this sends message m
with contents p to object y

Method Returns

▮ `return(v)` Current object this sends message return
with contents v to calling object



Communication History in Java

History of sequential Java prog: sequence of messages/events containing data

```
void main() {  
    Stack s = new Stack();    Stack t = new Stack();  
    s.push(5);    t.push(3); s.push(7);  
}
```

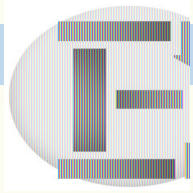
Global history of entire program:

(call_push(5) ret_push call_push(2) ret_push call_push(7) ret_push)

Local history of object s:

(call_push(5) ret_push call_push(7) ret_push)

With encapsulation: fully determines current state

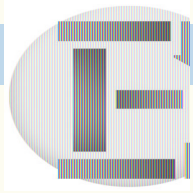


Why Histories?

- + Semantic foundation of OOP
 - + Fully Abstract Trace Semantics for a Core Java Language (Rathke & Jeffrey)
- + Behavioral Subtyping, fragile base class problem
 - + Behavioural Subtyping Relations for OO Formalisms (Fischer & Wehrheim)
- + Implementation independent
- + Goal I: finding a 'familiar' specification language for Java based on histories and amenable to automated run-time checking.

Goal II: allow *combining*

- + protocol-oriented properties (orderings of events) and data-oriented properties (assertion checking)



Example: BufferedReader

```
public class BufferedReader {  
    BufferedReader(Reader in)  
    BufferedReader(Reader in, int sz)  
  
    void    close()  
    void    mark(int readAheadLimit)  
    boolean markSupported()  
    int     read()  
    int     read(char[] cbuf, int off, int len)  
    String  readLine()  
    boolean ready()  
    void    reset()  
    long    skip(long n)
```



Communication View

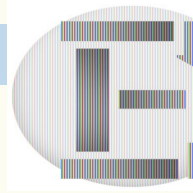
Partial mapping

From Event 'type'

To Name

```
local view BReaderHistory grammar BReader.g
specifies java.io.BufferedReader {
    new(Reader in)    OPEN,
    call String readLine()    READ,
    call void close()    CLOSE
}
```

- Specifies which events to capture (partial mapping)
- User-defined abstraction of local history (projection, identify >1 events with single terminal)



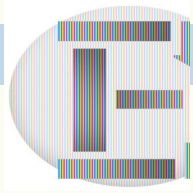
Attribute Grammars (I)

- Context-free Grammar + Attributes (map (sub)words of the language to values)

- Value of attributes defined in the productions

$$\begin{array}{lcl} S ::= S1 \ a & \{ & S.\text{height} = S1.\text{height} + 1; \} \\ & | \ a & \{ S.\text{height} = 0; \} \end{array}$$

- Synthesized Attr.: propagate info up the tree
 - Determines attribute val. of non-terminal in LHS
- Inherited Attr.: pass info down the tree
 - Determine attribute val. of non-terminals in RHS



Attribute Grammars (II)

Basic Idea:

1. Specification = set of (valid) histories = formal language generated by an (extended) attribute grammar
2. Run-time checking: parse trace in grammar + check asserts

Terminals in grammar:

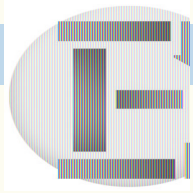
Message names: objects with built-in attr (caller, params, ...)

Non-terminals in grammar:

User-defined, specify protocol structure.

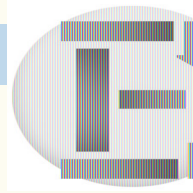
Attributes of non-terminals *define* data properties

Assertions *specify* properties of attributes



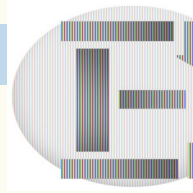
Attribute Grammar – BufferedReader client

```
start : OPEN  READ*  
      (CLOSE {assert OPEN.caller() == CLOSE.caller();})?  
      |  
       $\epsilon$ 
```

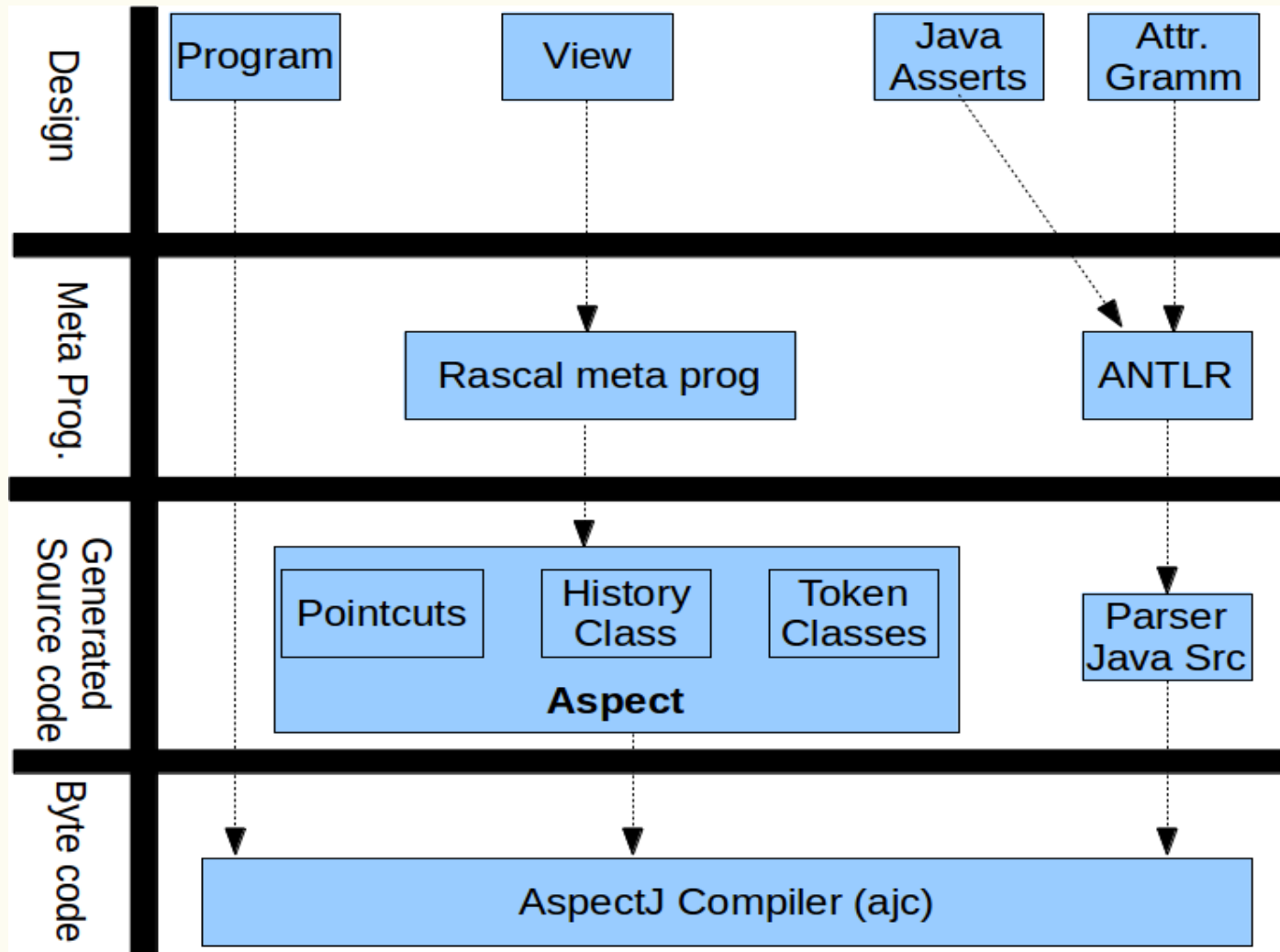


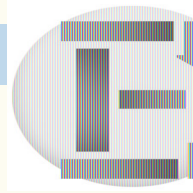
Why Attribute Grammars?

- Grammars are ‘familiar’: taught to undergraduates
- Allow **combining** protocol and data-oriented properties
- Declarative
- Amenable to automated run-time verification (well-developed parsing technology)



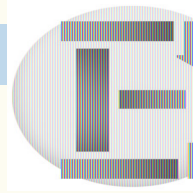
Tool Architecture





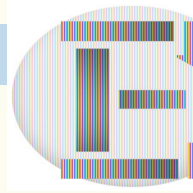
Challenges for industrial usage

- How to integrate checker into existing build process?
- Support for language features: Constructors, Inheritance / Dynamic Binding, Overloading, Static Methods, Multi-threading
- Instrumentation: source code not always available (Library methods/third party libs)
- Java Security model: classes to be monitored may not always be visible or accessible to monitor
- Writing specs is time consuming (refactorings?)



Non-functional Properties

- “Resource usage (mem/CPU/bandwidth) is above a certain threshold”
- Monitoring SLAs: “Average number of queries completed in the last 24h is ≥ 10 ”
- Currently: ad-hoc approach
- Manual reaction from operations team (domain-specific knowledge)
- Translation to functional properties (record time of event in attribute)? (interference of monitors)



Further Reading (selected publications)

Run-Time Assertion Checking of Data- and Protocol-Oriented Properties of Java Programs: An Industrial Case Study.

Frank S. de Boer, Stijn de Gouw, Einar Broch Johnsen, Andreas Kohn, Peter Y. H. Wong

[T. Aspect-Oriented Software Development 11: 1-26 \(2014\)](#)

Run-Time Verification of Coboxes.

Frank S. de Boer, Stijn de Gouw, Peter Y. H. Wong

[SEFM 2013: 259-273](#)

Run-Time Verification of Black-Box Components Using Behavioral Specifications: An Experience Report on Tool Development.

Frank S. de Boer, Stijn de Gouw

[FACS 2012:128-133](#)

<https://github.com/cwi-swath/saga>